# **Evolutionary Dynamics of Collective Action** in N-person Stag-Hunt Dilemmas

Jorge M. Pacheco<sup>1</sup>, Francisco C. Santos<sup>2</sup>, Max Souza<sup>3</sup>, Brian Skyrms<sup>4</sup>

- <sup>1</sup> *ATP*-group, CFTC & Departamento de Fisica da Universidade de Lisboa, Complexo Interdisciplinar, Av Prof. Gama Pinto 2, 1649-003 Lisboa, Portugal.
- <sup>2</sup> *IRIDIA/CoDE*, Université Libre de Bruxelles, Av. F. Roosevelt 50, CP 194/6, Brussels, Belgium,
- <sup>3</sup> Departamento de Matemática Aplicada, Universidade Federal Fluminense, R. Mário Santos Braga, s/n, Niterói-RJ, 24020-140, Brasil,
- <sup>4</sup> Logic and Philosophy of Science, School of Social Sciences, University of California at Irvine, Irvine, CA 92612, U.S.A.

**ELECTRONIC SUPPLEMENTARY MATERIAL** 

#### 1. N-PERSON STAG-HUNT IN INFINITE POPULATIONS

The evolutionary dynamics of Cs and Ds in the N-person Stag-Hunt game with a minimum threshold M can be studied by analyzing the sign of  $f_C - f_D$  (see Appendix 1). Hence, using the same conventions introduced in the Appendix 1, we shall study in detail the following polynomial

$$Q(x) = f_C - f_D = c \left(\frac{F}{N} - 1\right) - c \frac{F}{N} (1 - x)^{N - M} \sum_{k=0}^{M-1} {N-1 \choose k} (1 - M \delta_{k,M-1}) x^k (1 - x)^{M-1-k}.$$

The roots of Q(x) provide the interior fixed points of the replicator dynamics equation. In what follows, we shall assume that  $N \ge 2$ . For most of the time, we shall also assume that 1 < M < N. The degenerate cases will be dealt with at the end. Let us start by recasting Q(x) in a more amenable form. To this end, let  $F/N = \lambda$ ; we may rewrite

$$Q(x) = -c \left\{ 1 - \lambda + \lambda \left[ \sum_{k=0}^{M-1} {N-1 \choose k} x^k (1-x)^{N-1-k} - M {N-1 \choose M-1} x^{M-1} (1-x)^{N-M} \right] \right\}.$$

Since

$$1 = 1^{N-1} = (x+1-x)^{N-1} = \sum_{k=0}^{N-1} {N-1 \choose k} x^k (1-x)^{N-1-k},$$

we have that

$$Q(x) = -c \left\{ 1 - \lambda \left[ \sum_{k=M}^{N-1} \binom{N-1}{k} x^k (1-x)^{N-1-k} + M \binom{N-1}{M-1} x^{M-1} (1-x)^{N-M} \right] \right\}.$$

Let

$$R(x) = \sum_{k=M}^{N-1} {N-1 \choose k} x^k (1-x)^{N-1-k} + M {N-1 \choose M-1} x^{M-1} (1-x)^{N-M}$$

$$= x^{M-1} \left( \sum_{k=M}^{N-1} {N-1 \choose k} x^{k-M+1} (1-x)^{N-1-k} + M {N-1 \choose M-1} (1-x)^{N-M} \right)$$
(1)

Then we have that

$$Q(x) = -c(1 - \lambda R(x))$$

Hence, the roots of Q(x) are given by the intersection(s) of the line  $1/\lambda \equiv N/F$  with the polynomial R(x). It turns out that Figure 1-a provides examples of N/R(x), such that intersections with the line F identify the interior fixed points. We shall show below various properties of R(x) that capture the possibilities already illustrated in Figure 1, which we now prove are quite general.

## Lemma 1

- 1. R(0) = 0;
- 2. R(1) = 1;
- 3. R(x) > 0,  $x \in (0,1)$ ;
- 4. Let  $x^* = \frac{M}{N}$ . Then we have that R'(x) > 0 for  $0 \le x < x^*$ , and R'(x) < 0 for  $x^* < x < 1$ . In particular,  $R'(x^*) = 0$  and  $x^*$  is a point of maximum of R with  $R(x^*) > 1$ ;

Before we prove Lemma 1, let us use it to prove the main result:

# **Proposition 1**

Let  $\lambda^* = \frac{1}{R(x^*)}$ . We have that  $0 < \lambda^* < 1$ . Moreover, Q(x) satisfies:

- a. For  $\lambda < \lambda^*$  there are no roots in (0,1);
- b. For  $\lambda = \lambda^*$  there exists one double root at  $x = x^*$ ;
- c. For  $\lambda^* < \lambda \le 1$  there are two simple roots  $\{x_1, x_2\}$ , with  $x_1 \in (0, x^*)$  and  $x_2 \in (x^*, 1]$ ;
- d. For  $\lambda > 1$  there is a single root in  $(0, x^*)$ .

# **Proof of Proposition 1**

From Lemma 1 we have that  $R(x^*) > 1$ , thus  $0 < \lambda^* < 1$ . We then observe that

- i. For  $\lambda < \lambda^*$ , we have that  $\lambda R(x) < \lambda^* R(x^*) = 1$ . Thus Q(x) < -c(1-1) = 0
- ii. For  $\lambda = \lambda^*$ , we compute  $Q(x^*) = -c(1 \lambda^* R(x^*)) = -c(1 1) = 0$ . Also,  $Q'(x^*) = cR'(x^*) = 0$  and an easy calculation shows that  $R''(x^*) \neq 0$ . Hence,  $x^*$  is a double root.
- iii. For  $\lambda^* < \lambda \le 1$ , we first observe that we have Q(0) = -c,  $Q(1) = -c(1-\lambda) < 0$ . Since  $1 \lambda R(x^*) < 0$ , we have  $Q(x^*) > 0$ . By the Intermediate Value Theorem, Q(x) will have at least two roots: one in  $(0, x^*)$  and another at  $(x^*, 1]$ . Moreover, Q'(x) = cR'(x). Thus Q(x) is monotonically increasing in  $(0, x^*)$  and monotonically decreasing in  $(x^*, 0)$ . Thus these roots are unique.
- iv. For  $\lambda > \lambda^*$ , we now have Q(1) > 0, and thus there is no root in  $(x^*,1]$ . However, the argument for  $(0,x^*)$  remains unchanged, and we have the result. Let us now prove Lemma 1.

## **Proof of Lemma 1**

First, notice that (1), (2) and (3) are straightforward from the form of the polynomial R(x). cf. (Eq. 1). To prove (4), we let k = N - 1 - k', and given that

$$\binom{N-1}{N-1-k'} = \binom{N-1}{k'},$$

we may write

$$R(x) = x^{M-1} \left[ \sum_{k'=0}^{N-M-1} {N-1 \choose k'} x^{N-M-k'} (1-x)^{k'} + M {N-1 \choose M-1} (1-x)^{N-M} \right]$$
$$= x^{N-1} \left[ \sum_{k'=0}^{N-M-1} {N-1 \choose k'} \left( \frac{1-x}{x} \right)^{k'} + M {N-1 \choose M-1} \left( \frac{1-x}{x} \right)^{N-M} \right] .$$

Let  $z = \frac{1-x}{x}$ . Then, we have that  $z' = -\frac{1}{x^2} = -\frac{1}{x}(z+1)$ .

Thus

$$R(x) = x^{N-1}p(z), \quad p(z) = \sum_{i=0}^{N-M} a_i z^i,$$

where

$$a_i = \binom{N-1}{i}, \ 0 \le i < N-M \ \text{and} \ a_{N-M} = M \binom{N-1}{M-1}$$

We now compute R':

$$\begin{split} R'(x) &= (N-1)x^{N-2}p(z) - x^{N-2}p'(z)(z+1) \\ &= x^{N-2} \Big[ (N-1)p(z) - p'(z)(z+1) \Big] \\ &= x^{N-2} \Big[ (N-1)\sum_{i=0}^{N-M} a_i z^i - \sum_{i=1}^{N-M} i a_i z^i - \sum_{i=1}^{N-M} i a_i z^{i-1} \Big] \\ &= x^{N-2} \Big[ (N-1)a_0 - a_1 + (N-1)\sum_{i=1}^{N-M} a_i z^i - \sum_{i=1}^{N-M} i a_i z^i - \sum_{i=2}^{N-M} i a_i z^{i-1} \Big] \end{split}$$

Since  $a_0 = 1$  and  $a_1 = N - 1$ , and writing i = i + 1 in the last sum, we find that

$$R'(x) = x^{N-2} \left[ (N-1) \sum_{i=1}^{N-M} a_i z^i - \sum_{i=1}^{N-M} i a_i z^i - \sum_{i=2}^{N-M-1} (i+1) a_{i+1} z^i \right] = x^{N-2} S(z)$$

where

$$S(z) = \sum_{i=1}^{N-M-2} \left[ (N-1-i)a_i - (i+1)a_{i+1} \right] z^i + \left[ Ma_{N-M-1} - (N-M)a_{N-M} \right] z^{N-M-1} + (M-1)a_{N-M} z^{N-M} .$$

On noting that

$$\binom{L}{j+1} = \frac{L-j}{j+1} \binom{L}{j}$$
 (2)

we obtain, for  $1 \le i < N - N$ , that

$$a_{i+1} = \frac{N - 1 - i}{i + 1} a_i.$$

Hence,

$$\sum_{i=1}^{N-M-2} [(N-1-i)a_i - (i+1)a_{i+1}]z^i = 0 .$$

Also, we have

$$Ma_{N-M-1} - (N-M)a_{N-M} = M \binom{N-1}{M} - (N-M) \binom{N-1}{M-1},$$

which on calling upon (Eq. 2) yields

$$M \binom{N-1}{M} - (N-M) \binom{N-1}{M-1} = (N-M) \binom{N-1}{M} - (N-M) \binom{N-1}{M-1}$$
$$= -(N-M)(M-1) \binom{N-1}{M-1} .$$

Thus, we can write

$$S(z) = z^{N-M-1} \binom{N-1}{M-1} \left[ -(N-M)(M-1) + M(M-1)z \right]$$

which yields

$$R'(x) = x^{M-1} (1-x)^{N-M-1} {N-1 \choose M-1} [-(N-M)(M-1) + M(M-1)z]$$
 (3)

For  $x \in (0,1)$ , (Eq. 3) vanishes at

$$z^* = \frac{N-M}{M} = \frac{1-M/N}{M/N} \qquad .$$

Since 
$$z = \frac{1-x}{x}$$
,  $x^* = \frac{M}{N}$ .

Also, from (Eq. 3), we see that

- i. For  $0 < z < z^*$ , R'(x) < 0;
- ii. For  $z > z^*$ , R'(x) > 0.

Moreover,  $z = \frac{1-x}{x}$  is monotonically decreasing and maps (0,1) into  $(0,\infty)$  (thus reversing the orientation), which yields that  $0 < z < z^*$  corresponds to  $x^* < x < 1$  and  $x > z^*$  corresponds to  $x < x^*$ . This proves (4).

Next we consider the degenerate cases not included in the proofs above.

## Degenerate cases

For the cases, M = 1 and M = N the above analysis does not hold, but they can be easily analyzed directly. Since

$$p(z) = \sum_{i=0}^{N-M} {N-1 \choose i} z^{i} + (M-1) {N-1 \choose M-1} z^{N-M}$$

we have for M = 1 that

$$R(x) = x^{N-1}(z+1)^{N-1} = 1.$$

Thus  $Q(x) = -c(1 - \lambda)$ , with  $\lambda^* = 1$  and then Q(x) = 0.

For M = N, we have that

$$R(x) = Nx^{N-1}$$
 and  $Q(x) = -c(1 - \lambda Nx^{N-1})$ .

Thus Q will have a single root for  $\lambda > \lambda^* = 1/N$ .

In any case, for  $1 \le M \le N$ , we have that

$$R(x^*) = \left(\frac{M}{N}\right)^{N-1} \left[ \sum_{i=0}^{N-M} {N-1 \choose i} \left(\frac{N-M}{M}\right)^i + \left(M-1\right) {N-1 \choose M-1} \left(\frac{N-M}{M}\right)^{N-M} \right].$$

Recalling that  $\lambda^* = \frac{1}{R(x^*)}$  and that  $\lambda = \frac{F}{N}$ , we may write the critical F,  $F^*$ , as

$$F^* = N^N \left[ \sum_{i=0}^{N-M} \binom{N-1}{i} (N-M)^i M^{N-1-i} + (M-1) \binom{N-1}{M-1} (N-M)^{N-M} M^{-1} \right]^{-1}.$$

## 2. N-PERSON PRISONER'S DILEMMA IN FINITE POPULATIONS

Here we detail the derivation of  $f_C(k) - f_D(k)$  for the N-person Prisoner's Dilemma in finite, well-mixed populations. We may write

$$\begin{split} f_{C}(k) - f_{D}(k) &= \binom{Z-1}{N-1}^{-1} \sum_{j=0}^{N-1} \left\{ \binom{k-1}{j} \binom{Z-k}{N-j-1} \Pi_{C}(j+1) - \binom{k}{j} \binom{Z-1-k}{N-1-j} \Pi_{D}(j) \right\} = \\ &= c \binom{Z-1}{N-1}^{-1} \sum_{j=0}^{N-1} \left\{ \binom{k-1}{j} \binom{Z-k}{N-j-1} \frac{F}{N}(j+1) - \binom{k}{j} \binom{Z-1-k}{N-1-j} \frac{F}{N}j \right\} \end{split}$$

Introducing the notation  $\tilde{x} = x - 1$  we may now write

$$\begin{split} f_{C}(k) - f_{D}(k) &= c \binom{\widetilde{Z}}{\widetilde{N}}^{-1} \sum_{j=0}^{\widetilde{N}} \left\{ \binom{\widetilde{k}}{j} \binom{Z-k}{\widetilde{N}-j} \frac{F}{N} (j+1) - \binom{k}{j} \binom{\widetilde{Z}-k}{\widetilde{N}-j} \frac{F}{N} j \right\} = \\ &= c (\frac{F}{N}-1) + \frac{F}{N} \binom{Z-1}{N-1}^{-1} \sum_{j=0}^{\widetilde{N}} j \left\{ \binom{\widetilde{k}}{j} \binom{Z-k}{\widetilde{N}-j} - \binom{k}{j} \binom{\widetilde{Z}-k}{\widetilde{N}-j} \right\} \end{split}$$

We may readily simplify the complicated sum obtaining the desired result:

$$\begin{split} f_C(k) - f_D(k) &= c \left[ (\frac{F}{N} - 1) + \frac{F}{N} \left( \frac{\tilde{N}}{\tilde{Z}} \right) (\tilde{k} - k) \right] = c \left[ \frac{F}{N} \left( 1 - \frac{\tilde{N}}{\tilde{Z}} \right) - 1 \right] \\ &= c \left[ \frac{F}{N} \left( 1 - \frac{N - 1}{Z - 1} \right) - 1 \right] \,. \end{split}$$